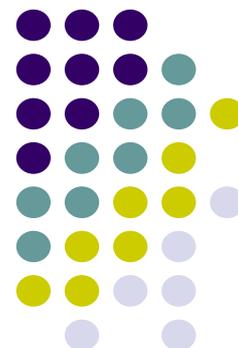


On Generalizations of Weighted Finite Automata and Graphics Applications

Jürgen Albert
Informatik II
University of Würzburg



Contents

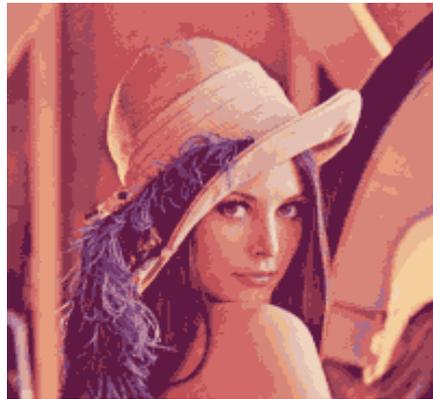
- Color, grayscale and bi-level images
- Color-cube
- Finite state acceptors, bi-level images
- Addressing in the unit-interval
- Bintreees
- Cartoons and video-clips
- Parametric WFA,
- Motivation, examples
- Iterated Function Systems
- Splines, texturing
- Open questions





Color Depths

24 Bits („true color“) vs. 4 Bits (color table) vs. 1 Bit

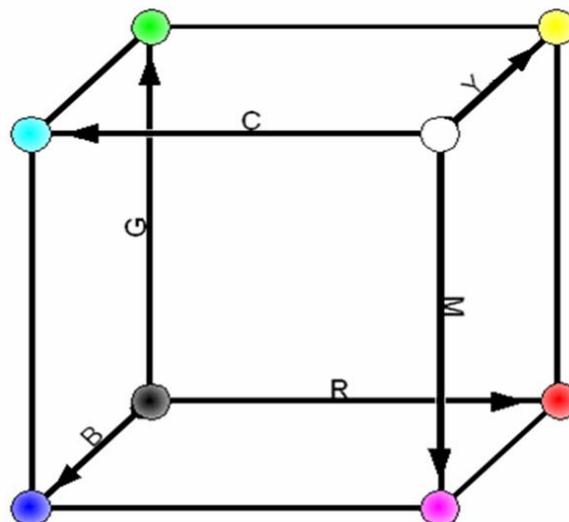


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Color-Cube

For custom digital photography in general 8 Bits for each component:
Black = (0, 0, 0), White = (255, 255, 255)



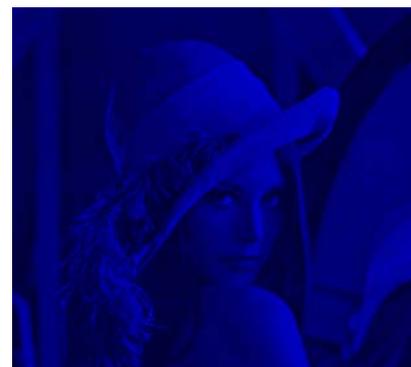
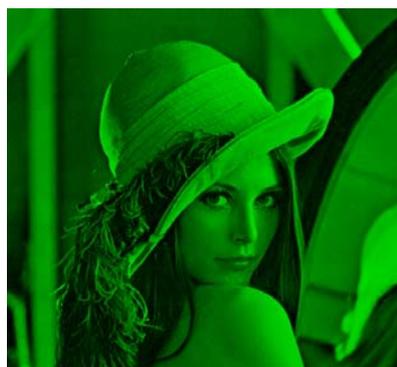
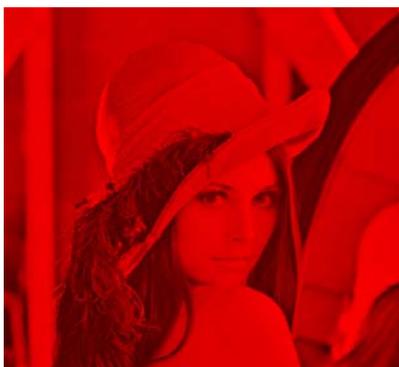
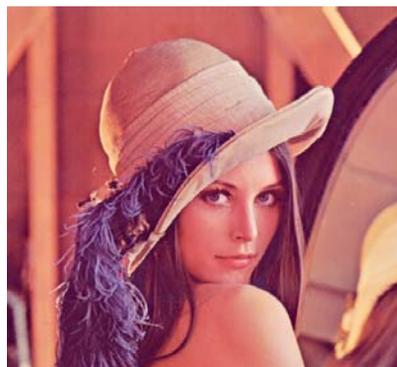
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Negatives - Complements to White



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Separation into Components (Red, Green, Blue)



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Colors + Global Transparency, „Alpha-Blending“ (1)



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Colors + Global Transparency, „Alpha-Blending“ (2)

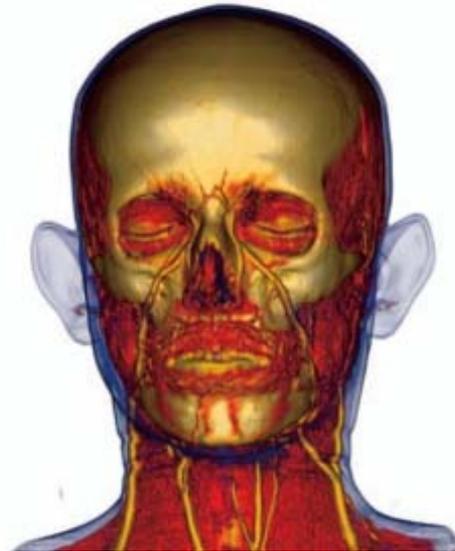


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Medical Imaging, Varying Transparency Values



Voxel densities often as 12-Bit-Integers, selecting density-intervals and rendering e.g. as in the „Visible Human Project“

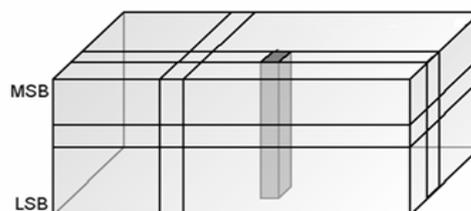
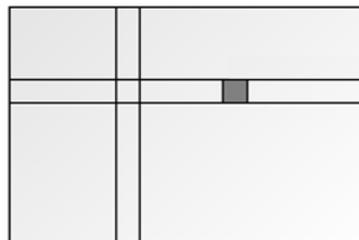


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Bitplane Encoding



- Each pixel-value is specified by its color components and its transparency via a fixed number of bits.
- Each of these bit-positions defines a bit-plane.
- The whole image is representable as a stack of bi-level images.



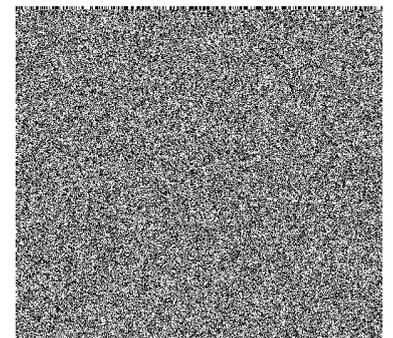
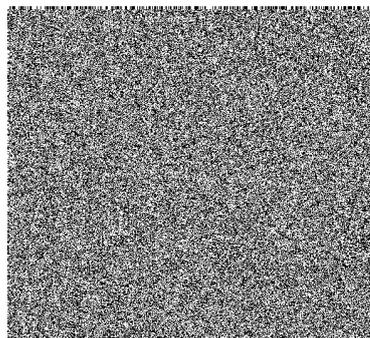
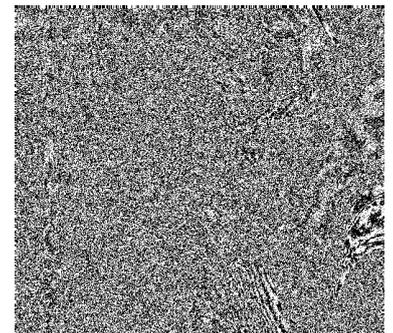
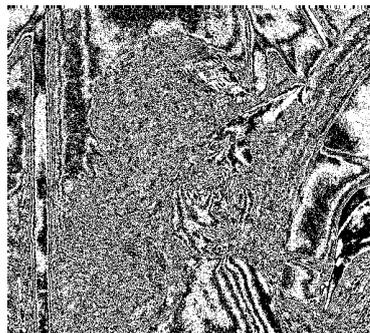
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8-Bit-Grayscale: Most Significant Bit-planes



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8-Bit-Grayscale: Least Significant Bit-planes



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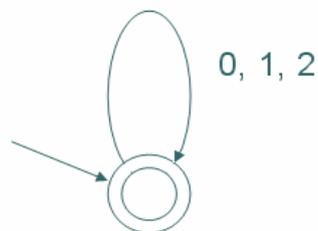
Words and Intervals

- For any alphabet $S = \{0, 1, \dots, p-1\}$
a word $w = a_1 a_2 \dots a_n$ in S^* is associated with the half-open interval $[\sum a_i p^{-i}, \sum a_i p^{-i} + p^{-n})$
- In the p -adic number-system this is denoted by $[0.a_1 a_2 \dots a_n, 0.a_1 a_2 \dots a_n + 0.00 \dots 1)$
- E.g. for $p = 4$ the word $w = 103$ addresses the half-open interval $[0.103, 0.110)$ of length 4^{-3}



Acceptors for Bi-level Images

Let $A = (\{q_1\}, \{0, 1, 2, 3\}, M, q_1)$ be a finite state acceptor with transitions



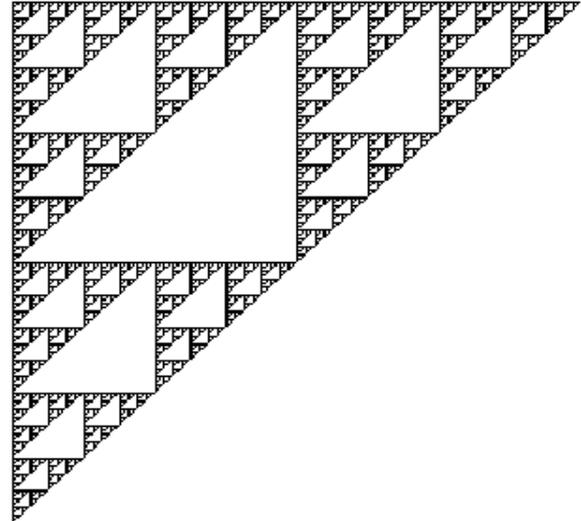
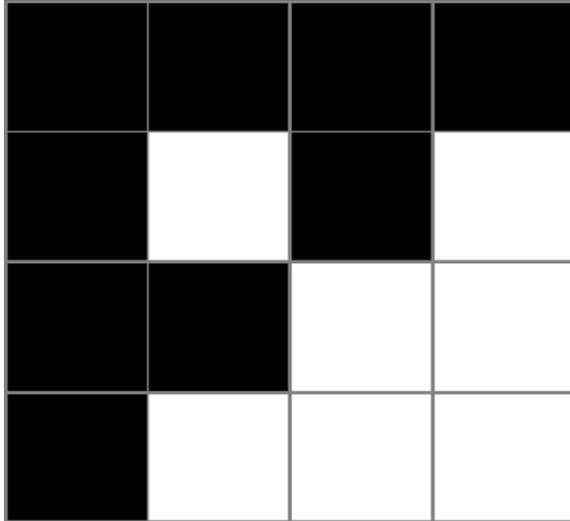
The acceptance pattern for sequences over $\{0, 1, 2, 3\}$ of length 2 is





2D-Interpretation

Acceptance patterns for automaton A and input-lengths 2 and 8.
Generation of bi-level images:

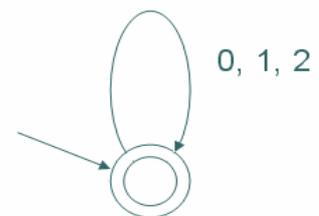
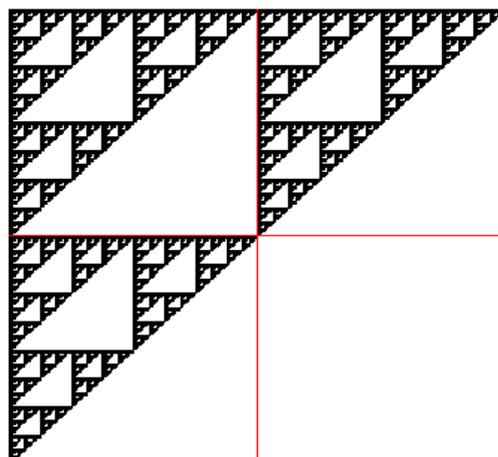


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Automaton Inference

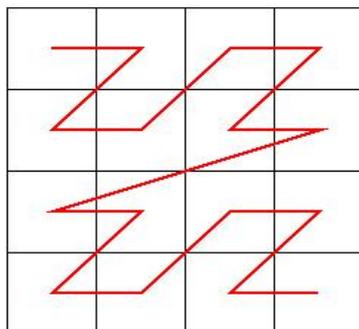
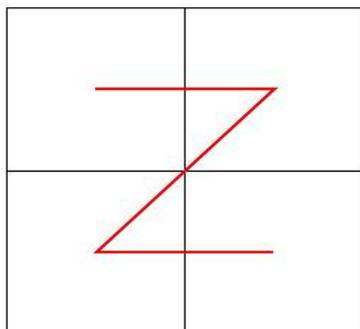
For simple cases a generating automaton can directly be inferred from the image. „Zooming“ into the quadrants 0, 1 and 2 yields scaled copies of the complete picture.



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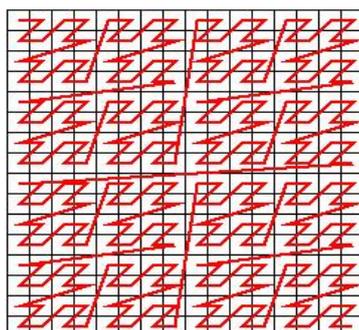
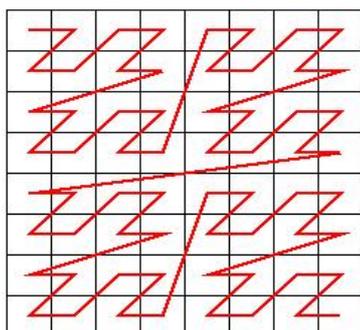


Morton-Order (Z-Order)



Addressing $2^k \times 2^k$ pixels in hierarchical and self-similar way.

Rotations and flips of Z are also very common in practice

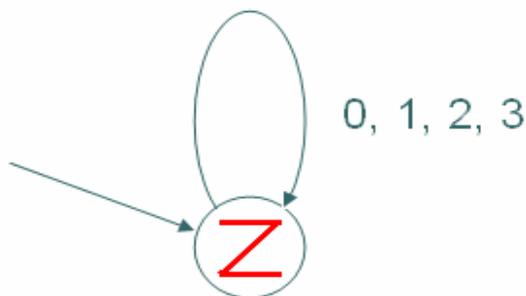


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Self-Similar Morton-Order

Self-similarity holds for the traversals of the Morton-Order in a trivial manner:

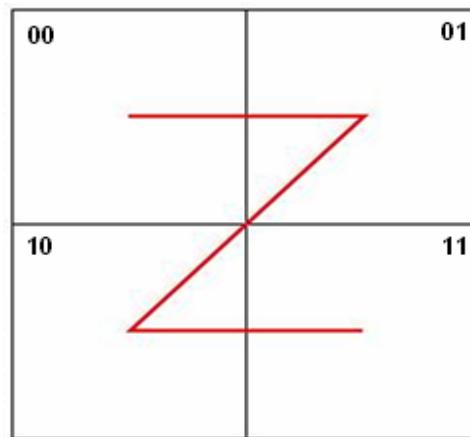


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Binary Alphabet

Addressing pixels / voxels requires only a binary alphabet „with interleaving of the dimensions“:

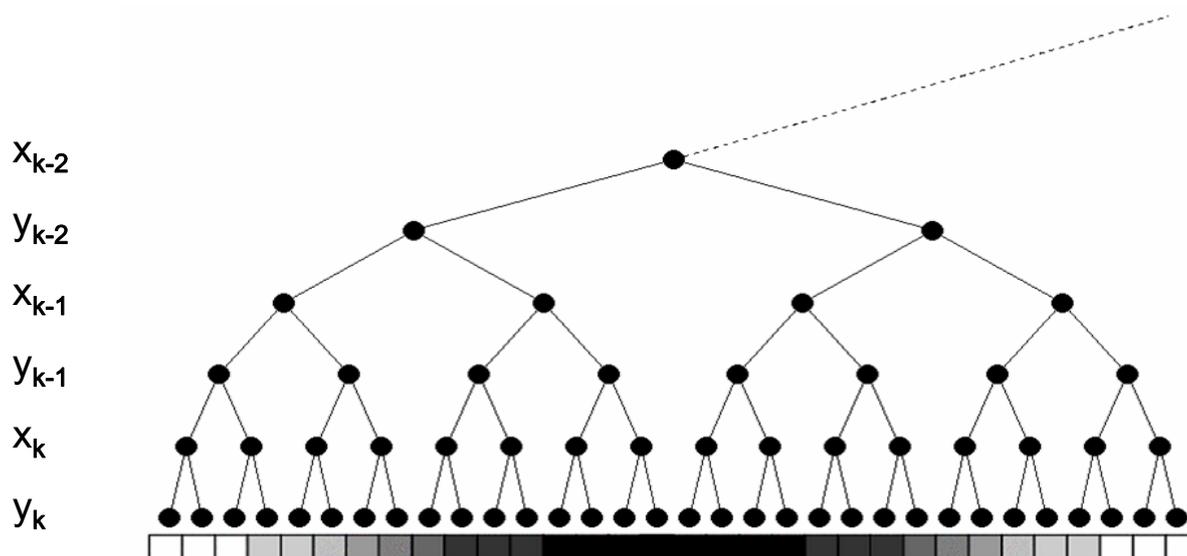


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2-D Bintree-Interpretation

Any image (black/white, grayscale, color) can be addressed by a bintree with interleaving the x- and y-bits of the addresses:

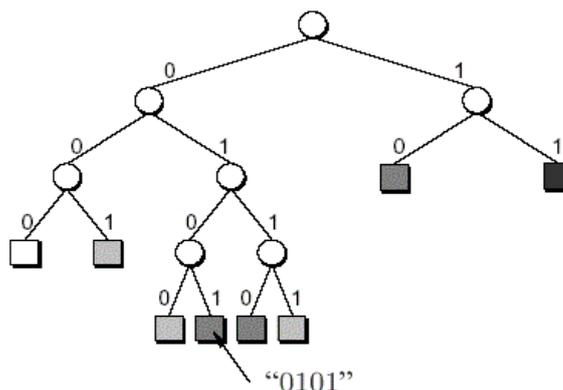
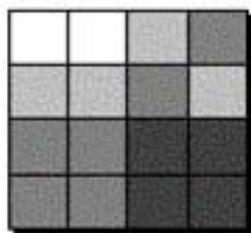


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Bintree Addressing

For any image, a tree can define a corresponding finite automaton: In its simplest form each pixel corresponds to an accepting state in a DFA, the root is the initial state.



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WFA - Definition

A **Weighted Finite Automaton (WFA)** $A = (Q, \Sigma, W_{\Sigma}, I, F)$ consists of:

- Set of **states** $Q = \{q_1, \dots, q_n\}$
- **Input-alphabet** $\Sigma = \{0, 1, \dots, p-1\}$, ($\{0, 1\}$ is sufficient)
- Set of **weight-matrices** $W_{\Sigma} = \{W_0, W_1, \dots\}$, where $W_i \in \mathbb{R}^{n \times n}$
- **Initial distribution** (row vector) $I \in \mathbb{R}^n$
- **Final distribution** (column vector) $F \in \mathbb{R}^n$

The image attached to the input-sequence $w = i_1 i_2 \dots i_t$ („the address“) is computed by A as:

$$f_A(w) = I \cdot W_{i_1} \cdot W_{i_2} \cdot \dots \cdot W_{i_t} \cdot F$$

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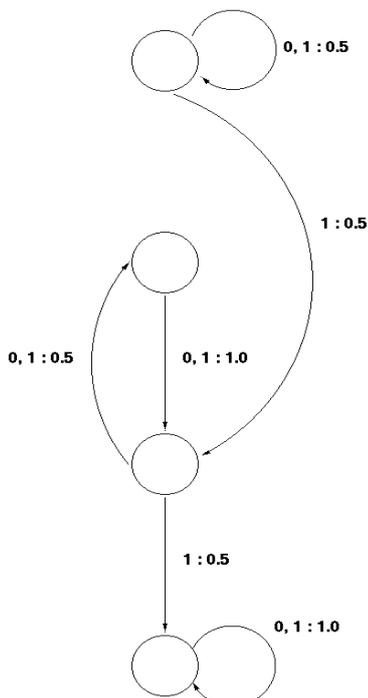
WFA - Example

Set of states $Q = \{q_1, q_2, q_3, q_4\}$
Set of input-symbols $\Sigma = \{0, 1\}$
Transition-matrices $W = \{M_0, M_1\}$
Initial distribution $I = (0, 0, 0, 1)$
Final distribution $F = (1, 0.5, 0.5, 0.25)^T$

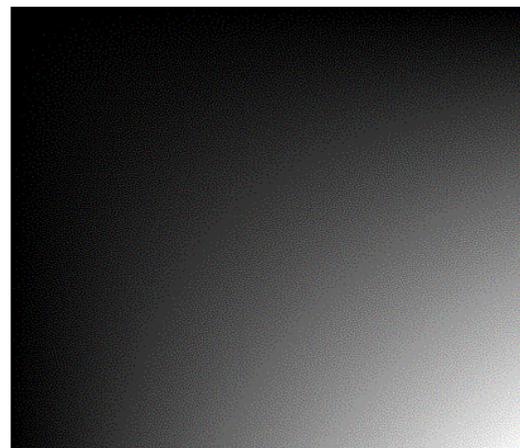
$$M_0 = \begin{pmatrix} 1.0 & 0.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.5 & 0.0 \\ 0.0 & 1.0 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.5 \end{pmatrix}, M_1 = \begin{pmatrix} 1.0 & 0.0 & 0.0 & 0.0 \\ 0.5 & 0.0 & 0.5 & 0.0 \\ 0.0 & 1.0 & 0.0 & 0.0 \\ 0.0 & 0.5 & 0.0 & 0.5 \end{pmatrix}$$



WFA - Example



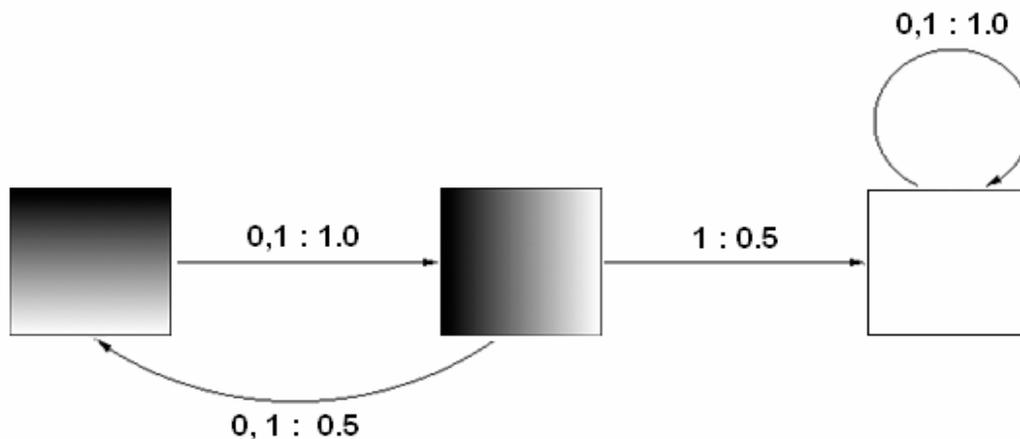
Graph for weight-matrices and grayscale image of the computed function $x \bullet y$ for $x, y \in [0, 1)$





WFA – Example, Subautomata

The WFA for $x \bullet y$ contains the subautomata for the linear slopes in y and x and the constant 1.



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WFA – Polynomials

$p_m(x) = x^m$, $m \geq 1$. Let $b_i \in \{0,1\}$ and

$x_1 = 0.b_1 b_2 b_3 \dots b_t$, $x_2 = 0.b_2 b_3 \dots b_t$,

$$\begin{aligned}
 x_1^m &= \left(\frac{1}{2}b_1 + \frac{1}{2}x_2 \right)^m \\
 &= \frac{1}{2^m} \sum_{i=0}^m \binom{m}{i} b_1^{m-i} \cdot x_2^i.
 \end{aligned}$$

For $b_1 = 0$:
$$x_1^m = \frac{1}{2^m} x_2^m$$

For $b_1 = 1$:
$$x_1^m = \frac{\binom{m}{0}}{2^m} x_2^0 + \frac{\binom{m}{1}}{2^m} x_2^1 + \dots + \frac{\binom{m}{m}}{2^m} x_2^m.$$

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WFA - Polynomials

- WFA can represent polynomials in a very economical way.
- Any polynomial of degree m can be computed by some WFA with $m+1$ states ("Line-Automata")
- Polynomials are the only smooth functions that WFA can generate exactly for arbitrary resolution. (Culik, Karhumäki, Kari, Steinby, Droste)
- Even the square-root function can only be approximated well, (Karhumäki, Terlutte, et. al.)

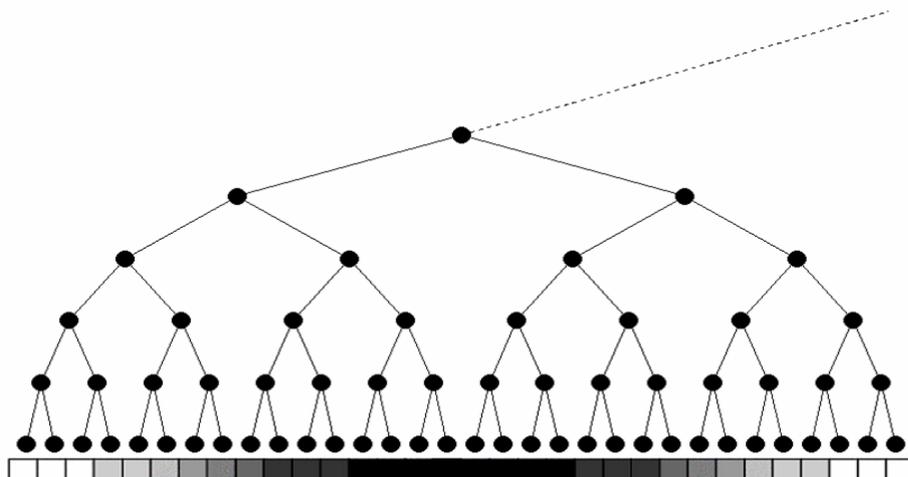
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WFA Inference Problem

For a given function (a k -dimensional grid of real-valued intensities) find a WFA which **approximates this function well**, and which can be stored in a **small number of bytes**.

Efficient heuristic implemented by Culik and Kari.



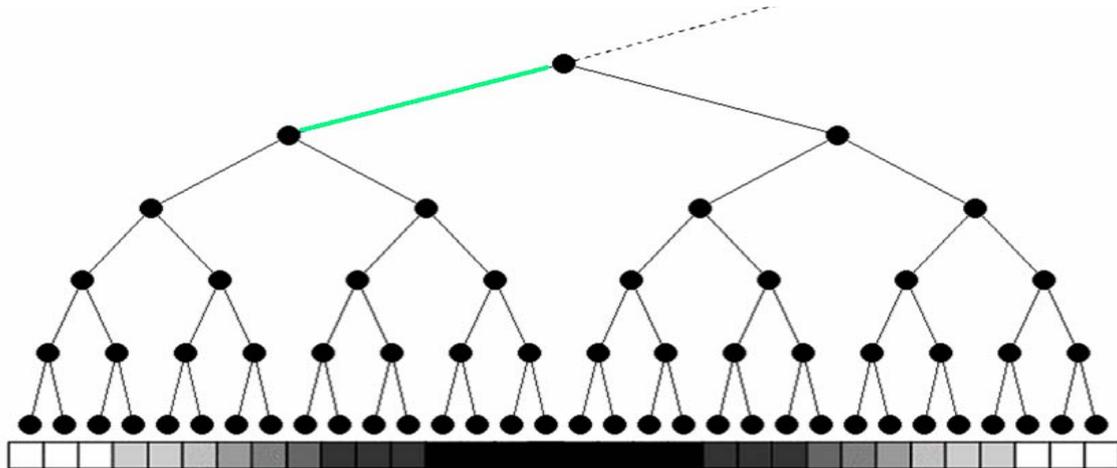
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Coding Decisions in the Bintree

At each inner node recursively check, whether to

- approximate the current sub-image by a **linear combination** or
- **subdivide** it further

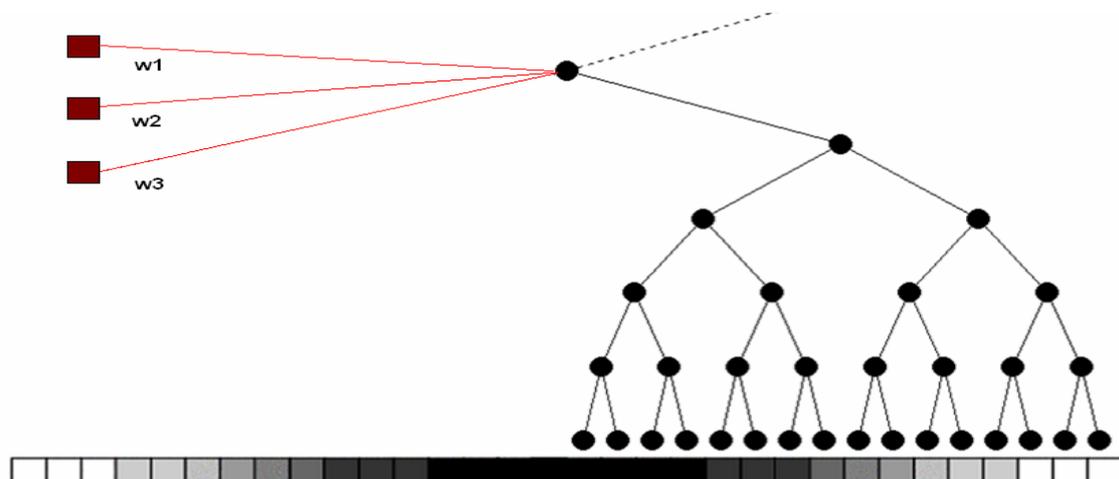


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Coding Decisions in the Bintree

Choosing a linear combination:

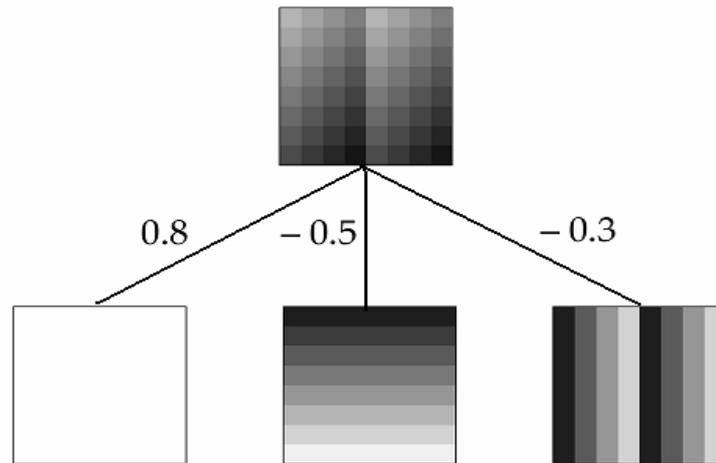


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Cost-Function

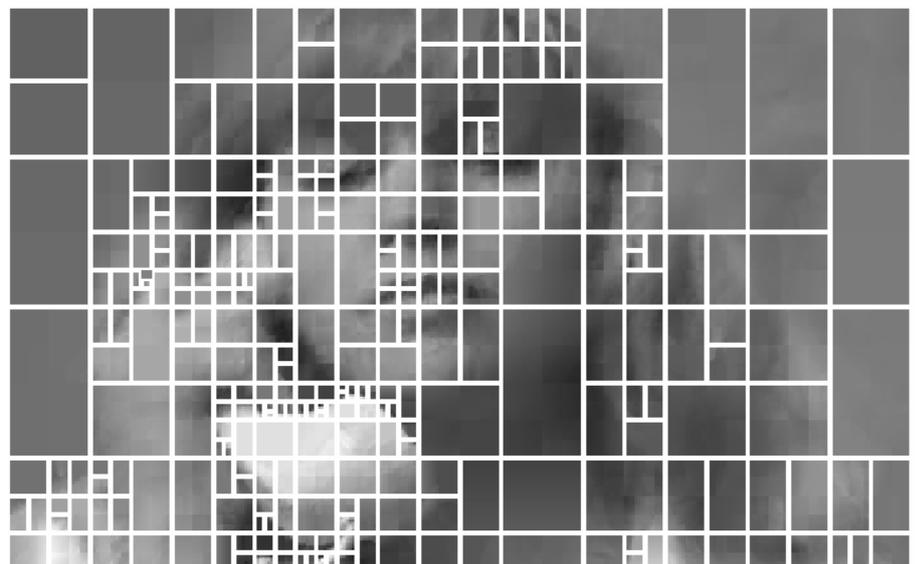
A cost-function compares the locally generated **error** and the storage **space** needed for the current approximation and thus checks the build-up of linear combinations in each step



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Expl. Bintree Decomposition



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JPEG Baseline vs. WFA for Low-Bitrates



JPEG 2000 vs. WFA



Testimage „Cafe“, (detail) 390 x 280 x 8bpp:

original,

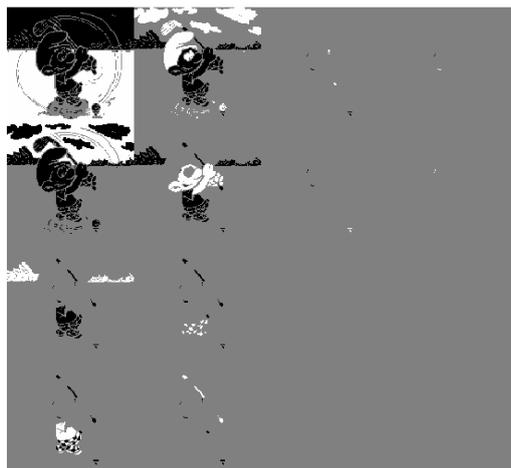
JPEG 2000,

WFA





WFA for Cartoon-like Images



- Separate all colors into bit-planes.
- Reconstruct by painting **layer over layer**.
- „**Don't care**“ values (in gray) are those which will be repainted later.
- Replace addition and multiplication in linear combinations of WFAs by **Boolean operators**. In experiments **XOR** showed best performance.
- In the cost-function count the number of wrong pixels to control error-rate vs. file-size.

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WFAs and Videosequences

Exploiting temporal redundancies similar to MPEG-x,
Different characteristics of frame-sequences:



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Motion Estimation in Fast Actions, „Head and Shoulder“ Sequences



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MWFA-Image-Partitioning and Motion-Vectors



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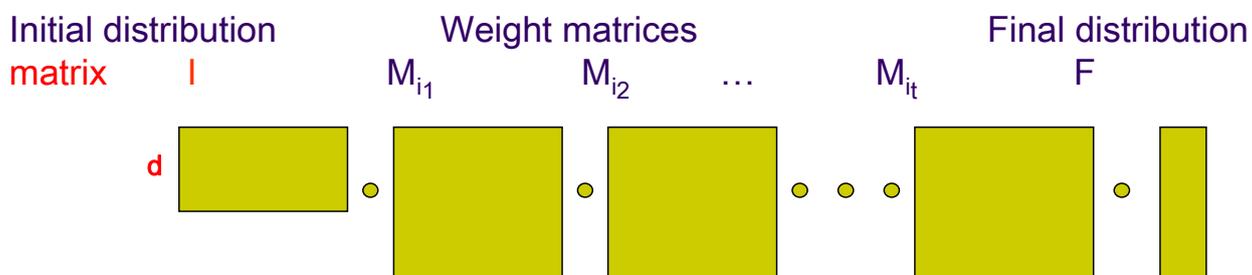
Motivations for Extending WFA

- In many practical applications there are finite higher dimensional grids for space, time, color-space, ...
- To display e.g. a square-root function is not really harder than to draw a square-function on the screen
- There are other types of self-similarities not covered by WFA



Parametric WFA (PWFA)

PWFA generalize WFA by replacing the **initial distribution vector** by a **matrix**. Now a **d-dimensional vector of real values** is computed for arbitrary input sequences $w = i_1 i_2 \dots i_t$.



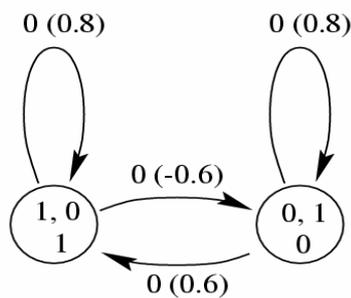
A point x belongs to the result-set $S(A)$ defined by a PWFA A , if there are infinitely many w s.t. $x = f_A(w)$ or the are points $f_A(w)$ arbitrarily close to x . The resulting d -dimensional vectors x are interpreted as **relations**, e.g. for $d=3$: bi-level in 3D or grayscale in 2D



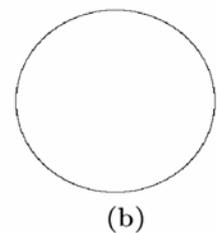
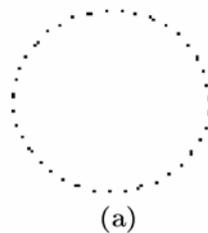
PWFA over Unary Alphabet

Rotations by an irrational angle e.g. $\cos^{-1}(0.8)$ applied to the final distribution point $(1, 0)$ yield the points of the unit circle.

$$S(A) = \{(\cos(t), \sin(t)) \mid t \in \mathbb{R}\}$$

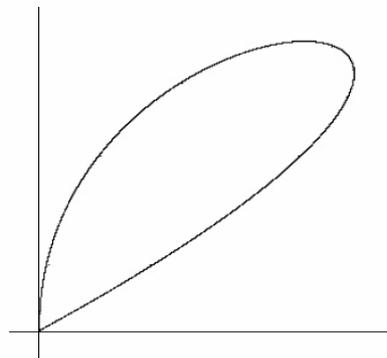


First 50 and 1000 points:



Polynomial Curves

For $(t^2 - t^3, t - t^2)$, $0 \leq t \leq 1$, a PWFA with 4 states can compute the 2D-curve-segment:

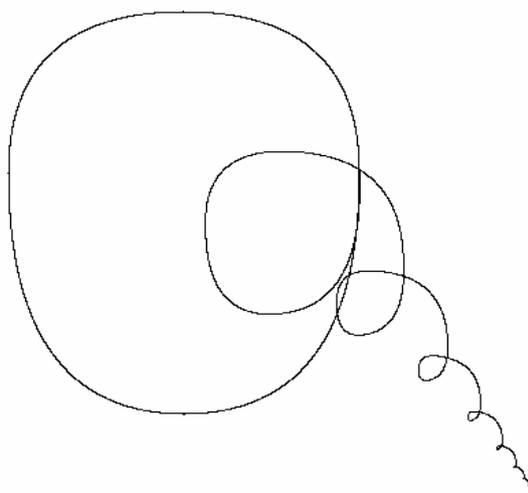


Any d -dimensional curve with parametric representation by d polynomials of maximal degree m can be computed by a PWFA with $m+1$ states.



Segments and Splines

Expl.: Piecewise combination of parabola-chunks
(9 states PWFA):

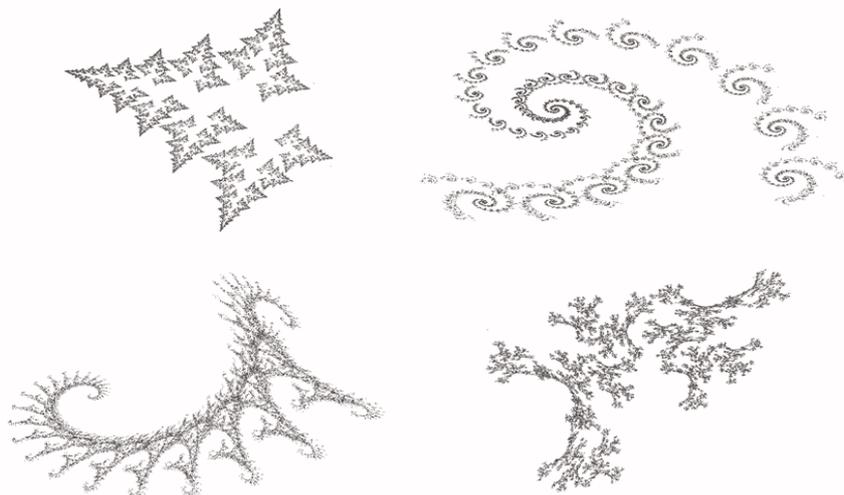


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Fractal Compression

Exploiting self-similarities in the pictures is comparable to dictionary-methods in text-compression in that „a new part of a picture is described by one or more references to previously coded parts“.



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Iterated Function Systems (IFS)

Affine Transformations in \mathbb{R}^m form the basis for IFS.

In \mathbb{R}^2 these are just:

$$w \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} + \begin{pmatrix} e \\ f \end{pmatrix} = A\bar{x} + B$$

Thus, affine transformations can contain scaling, rotation and translation.

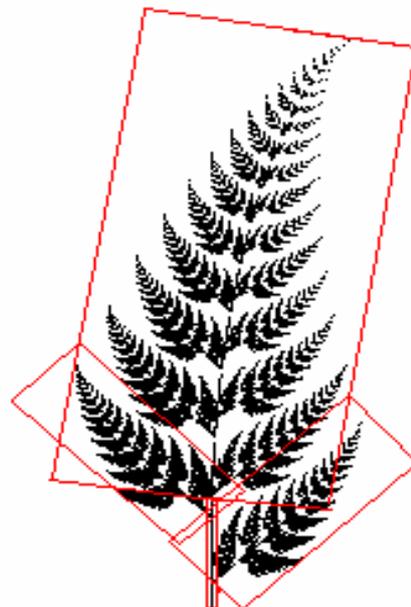
A transformation $f : X \rightarrow X$ is „contractive“, if there is an s such that $0 \leq s < 1$ and $d(f(x), f(y)) \leq s d(x, y)$ for all x, y in X .

IFS are defined via sets of contractive affine transformations.



Affine Transformations for „Fern“

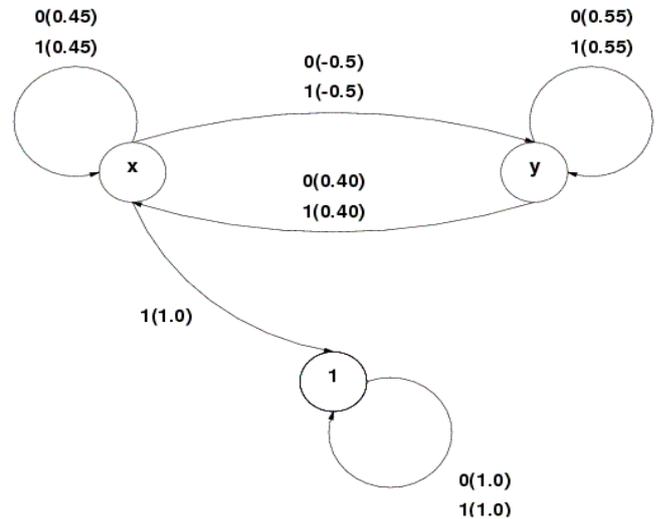
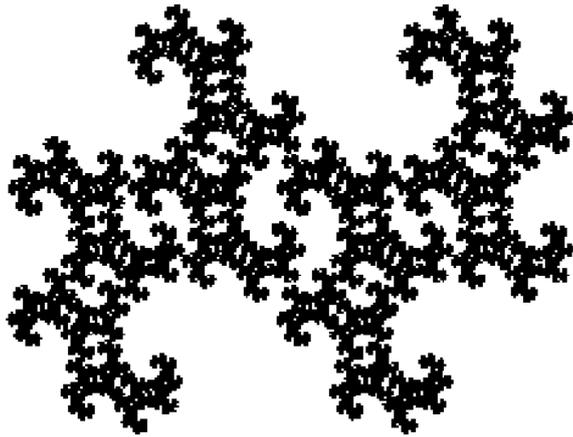
Four affine transformations for: upper part, new left and right leaf, new part of stem



PWFA-Simulation of Iterated Function Systems



Any 2D IFS with k contractive affine maps can be simulated by a PWFA with 3 states and k labels. Expl. „Dragon“:

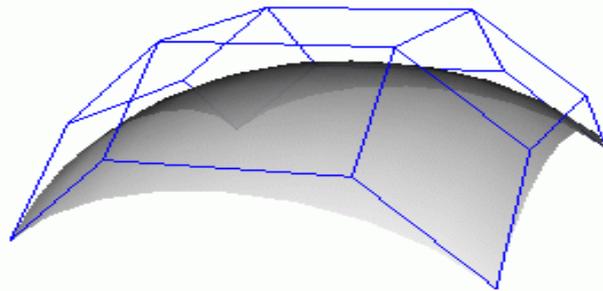


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Spline Surfaces



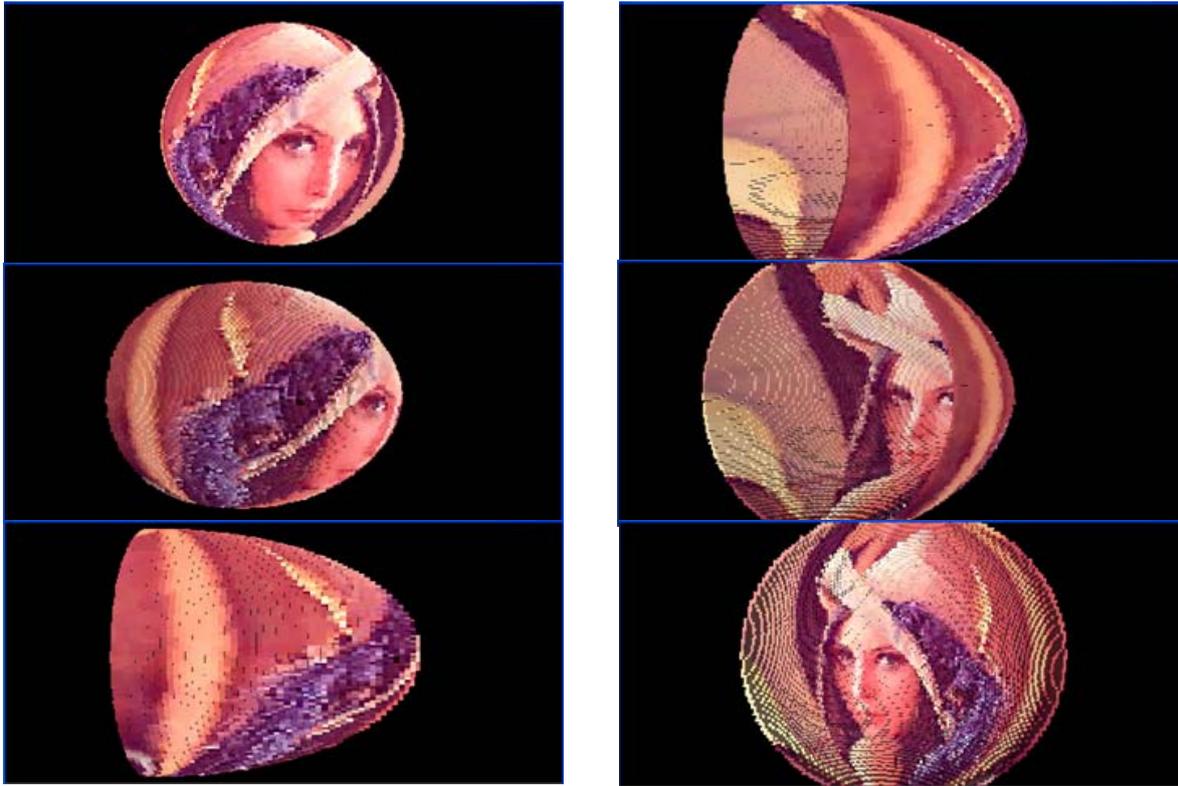
- E.g. the Bezier spline surfaces (or patches) are constructed of two Bezier curves
- The control points now make a control polyhedron of the surface.
- Moving the control points of one Bezier curve along a set of Bezier curves to define a surface



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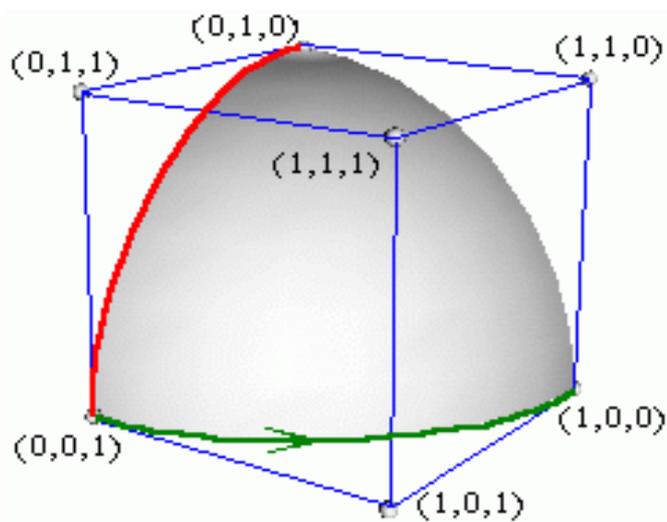
PWFA Spline Patches and Textures



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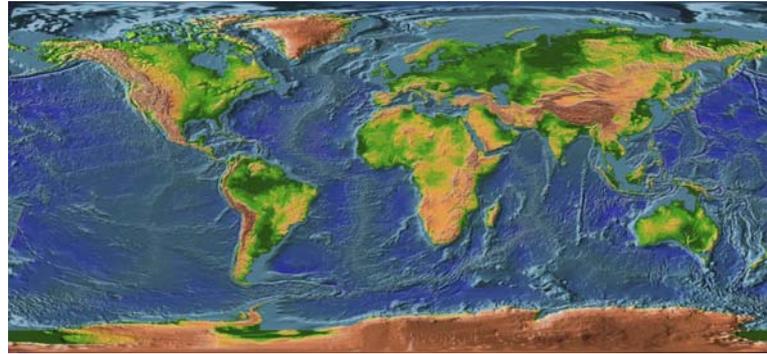
Spheres and Control Points



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Spheres and Textures



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Open Problems

- Pure WFA-encoding of video-clips and volume-data
- WFA: Encoding efficiency: „space for time“
- Inference heuristics for PWFA or for PWFA-subfamilies
- Applications of PWFA in „augmented reality“
- Efficiency for the representation of 3D-spline-patches
-

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